

LOCAL MONOTONICITY OF SUBGRADIENT MAPPINGS

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Global Monotonicity of Set-Valued Mappings

\mathcal{H} = Hilbert space, $T : \mathcal{H} \rightrightarrows \mathcal{H}$ set-valued mapping, operator

$$\text{gph } T = \{(x, y) \mid y \in T(x)\}, \quad \text{dom } T = \{x \mid T(x) \neq \emptyset\}$$

Monotonicity: $y_0 \in T(x_0), y_1 \in T(x_1) \implies \langle x_1 - x_0, y_1 - y_0 \rangle \geq 0$

Maximality: \nexists monotone T' with $\text{gph } T'$ strictly \supset $\text{gph } T$

Basic problem to solve for a given T mapping

Find x such that $0 \in T(x)$

Resolvants: $(I + cT)^{-1}, c > 0$ nonexpansive from monotonicity

Proximal point algorithm: $x^{k+1} = (I + c_k T)^{-1}(x^k), c_k \nearrow c_\infty$
under max monotonicity, $\{x^k\}$ converges to a solution (if \exists)

Subgradients Beyond Convex Analysis

function $f : \mathcal{H} \rightarrow (-\infty, \infty]$, lower semicontinuous, $\neq \infty$

Regular (“Frêchet”) subgradients: $y \in \widehat{\partial}f(x)$
 $f(x') \geq f(x) + \langle y, x' - x \rangle + o(\|x' - x\|)$

General (“limiting”) subgradients: $y \in \partial f(x)$
 $\exists \{x^\nu\}, \{y^\nu\}$, with $y^\nu \in \widehat{\partial}f(x^\nu)$, such that
 $y^\nu \xrightarrow{w} y, \quad x^\nu \rightarrow x, \quad f(x^\nu) \rightarrow f(x) < \infty$

Monotonicity characterization of the mapping $\partial f : \mathcal{H} \rightrightarrows \mathcal{H}$

∂f max monotone $\iff \partial f$ monotone $\iff f$ convex

Proximal point algorithm: for solving $0 \in \partial f(\bar{x})$ (optimization)

$$x^{k+1} = (I + c_k \partial f)^{-1}(x^k) \iff x^{k+1} = \operatorname{argmin}_x f^k$$

for the function $f^k(x) := f(x) + \frac{1}{2c_k} \|x - x^k\|^2$

application to dual problem underlies ALM in convex programming

Localization for Nonconvex Functions?

Local monotonicity: of $T : \mathcal{H} \rightrightarrows \mathcal{H}$ at $(x^*, y^*) \in \text{gph } T$
the monotonicity holds in neighborhood $\mathcal{X} \times \mathcal{Y}$ of (x^*, y^*)
likewise defined: local max monotonicity

Localized algorithm: for finding \bar{x} with $0 \in T(\bar{x})$
proximal point iterations succeed if started near a solution!

Local minimization: $0 \in \partial f(\bar{x})$ is necessary for local optimality
find such \bar{x} by local execution of the proximal point algorithm?
 $x^{k+1} = \text{a "subcritical point" of } f^k: 0 \in \partial f^k(x^{k+1})$

Challenges

- articulating the algorithm with local minimization steps
- implications for f of merely local monotonicity of ∂f ?????
- trouble from the weak convergence in the definition of ∂f

\implies henceforth assume the Hilbert space \mathcal{H} is finite-dimensional

Local Monotonicity From Variational Convexity

function $f : \mathcal{H} \rightarrow (-\infty, \infty]$, lower semicontinuous, $\neq \infty$
for simplicity, let $f(x)$ depend continuously on $(x, y) \in \text{gph } \partial f$

Variational convexity: of f at x^* for $y^* \in \partial f(x^*)$

\exists open convex neighborhood $\mathcal{X} \times \mathcal{Y}$ of (x^*, y^*) and some
convex lsc function \hat{f} on \mathcal{X} such that, for $(x, y) \in \mathcal{X} \times \mathcal{Y}$,
 $(x, y) \in \text{gph } \partial f \iff (x, y) \in \text{gph } \hat{\partial} \hat{f} \implies f(x) = \hat{f}(x)$

Consequences of variational convexity

- (a) the mapping $T = \partial f$ is maximal monotone in $\mathcal{X} \times \mathcal{Y}$
- (b) $f(x') \geq f(x) + \langle y, x' - x \rangle$ for all $x' \in \mathcal{X}$,
when $(x, y) \in \text{gph } \partial f$ in $\mathcal{X} \times \mathcal{Y}$

moreover **equivalence** holds when $y^* \in \hat{\partial} f(x^*)$

Strong Monotonicity and Variational Strong Convexity

Strong monotonicity: of $T : \mathcal{H} \rightrightarrows \mathcal{H}$ in $\mathcal{X} \times \mathcal{Y}$ at level $s > 0$
 $\langle x_1 - x_0, y_1 - y_0 \rangle \geq s \|x_1 - x_0\|^2$ for all $(x_i, y_i) \in [\mathcal{X} \times \mathcal{Y}] \cap \text{gph } T$
maximal: if $\text{gph } T$ can't be enlarged in $\mathcal{X} \times \mathcal{Y}$ with this maintained

Variational strong convexity: of f at level $s > 0$
same as earlier except with \hat{f} strongly convex at level s on \mathcal{X}
this corresponds to $f(x) - \frac{s}{2} \|x\|^2$ being convex on \mathcal{X}

Consequences of variational strong convexity

- (a) $T = \partial f$ is strongly maximal monotone at level s in $X \times Y$
- (b) $f(x') \geq f(x) + \langle y, x' - x \rangle + \frac{s}{2} \|x' - x\|^2$ for all $x' \in X$,
when $(x, y) \in \text{gph } \partial f$ in $X \times Y$

moreover **equivalence** holds when $y^* \in \hat{\partial} f(x^*)$

Variational Convexity Versus Convexity

Local convexity from variational convexity?

variational convexity of f for $\mathcal{X} \times \mathcal{Y} \implies$ convexity of f on \mathcal{X}
when $\mathcal{Y} = \mathcal{H}$ or when $f = f_0 + \delta_C$ for C closed convex, $f \in \mathcal{C}^1$

Example 1: $f(x) = \max\{1 - e^x, 1 - e^{-x}\}$ for $x \in \mathbb{R}$
graph = two “concave wings” rising from the origin of $\mathbb{R} \times \mathbb{R}$

f is variationally convex at $x^* = 0$ for $y^* = 0$

Example 2: $f(x_1, x_2) = x_2 + \delta_P(x_1, x_2)$, $P = \{(x_1, x_2) \mid x_2 = x_1^2\}$
graph = a “tilted parabola” floating in $\mathbb{R} \times \mathbb{R} \times \mathbb{R}$

f is variationally strongly convex everywhere, yet $\text{dom } f$ nonconvex

References

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downloads: sites.washington.edu/~rtr/mypage.html